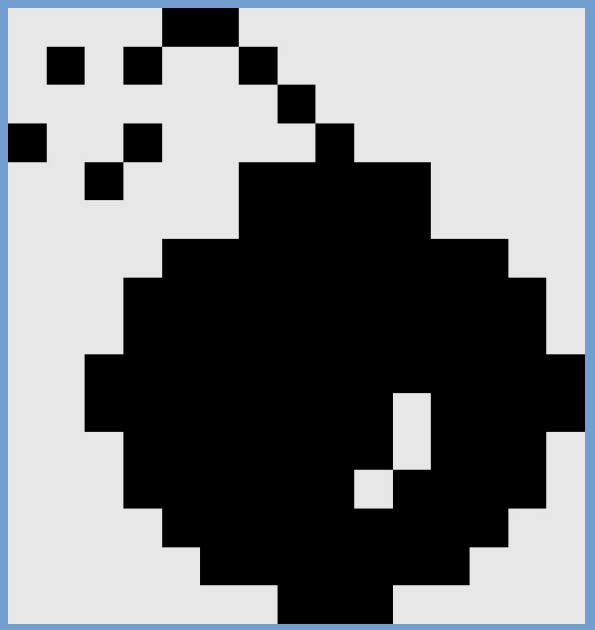
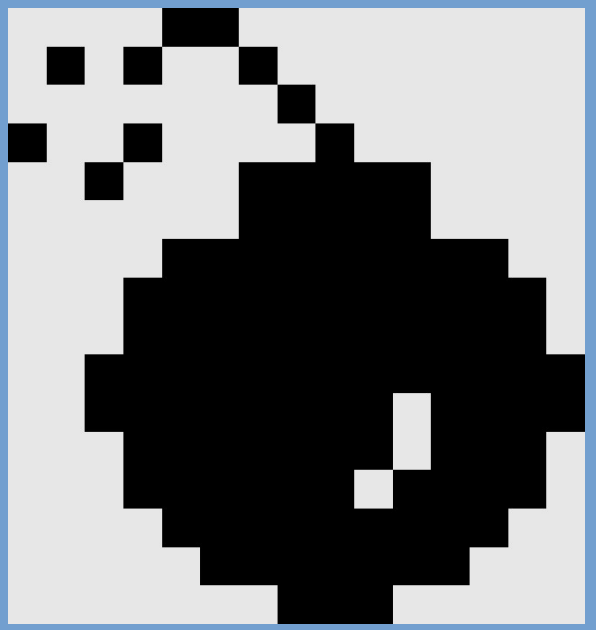
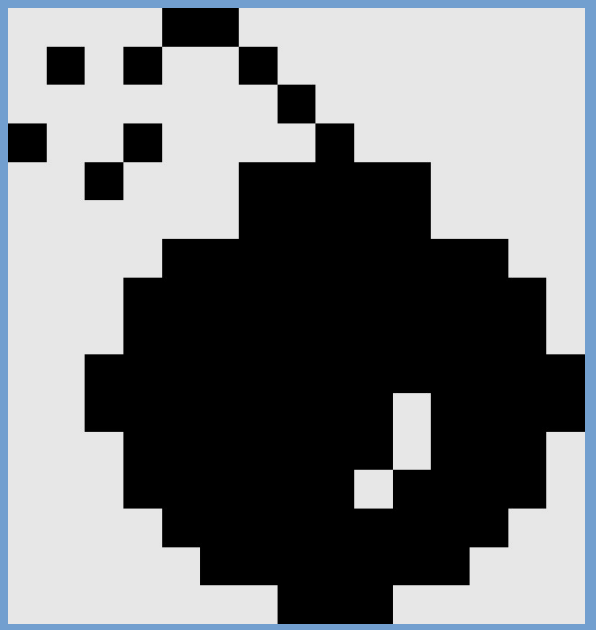

Fine-Grained Control-Flow Integrity through Binary Hardening

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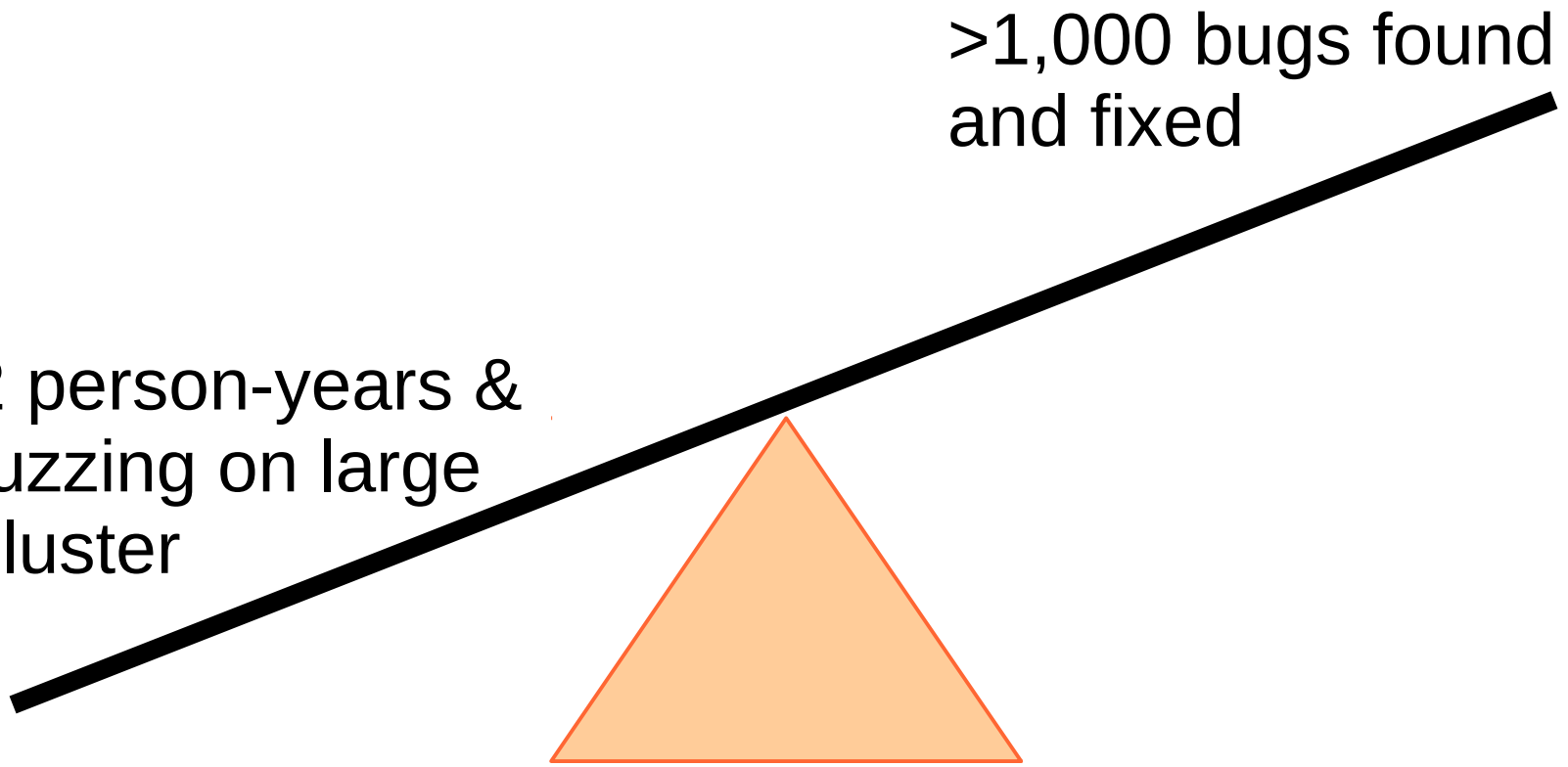
ETH Zürich

FFmpeg and a thousand fixes



http://j00ru.vexillium.org/?p=2211 Jan-10, 2014

Software is unsafe and insecure

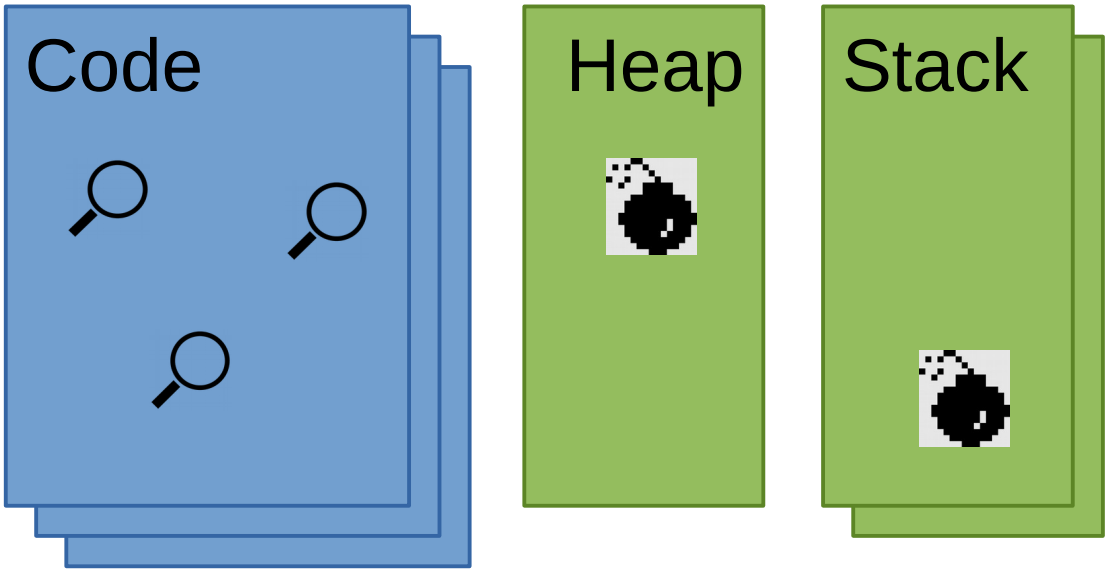
- Low-level languages (C/C++) trade type safety and memory safety for performance
  - Programmer responsible for all checks
- Large set of legacy and new applications written in C / C++ prone to memory bugs
- Too many bugs to find and fix manually
  - Protect integrity through safe runtime system


Vulnerability
Just Ahead

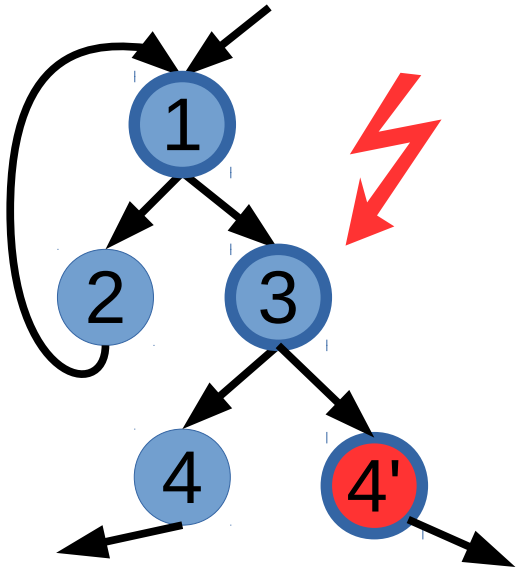
Code Reuse Attacks

Attack scenario: code reuse

- Find addresses of gadgets
- Force memory corruption to set up attack
- Leverage gadgets for code-reuse attack



Control-flow hijack attack



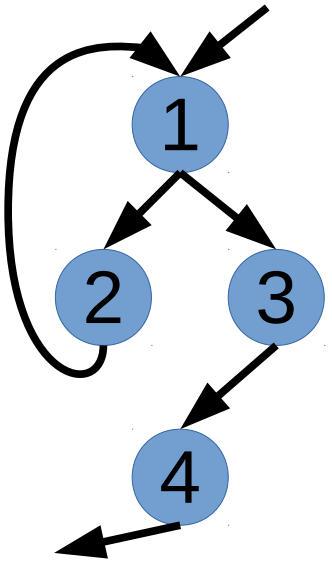
- Attacker modifies ***code pointer***
  - Function return
  - Indirect jump
  - Indirect call
- Control-flow leaves ***valid graph***
- Reuse existing code
  - Return-oriented programming
  - Jump-oriented programming

Control-Flow Integrity

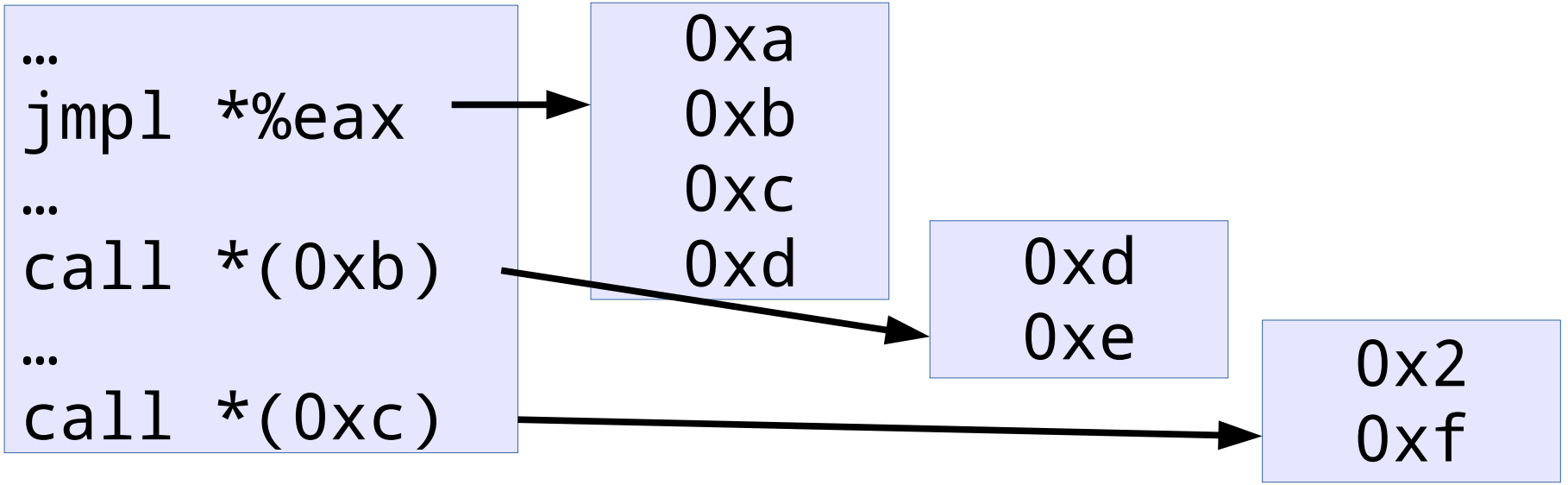
Control-Flow Integrity (CFI)

- CFI enforces that each dynamic indirect control flow transfer must target a statically determined set of locations
- Three sources of indirect transfers
  - Indirect jump
  - Indirect call
  - Function returns

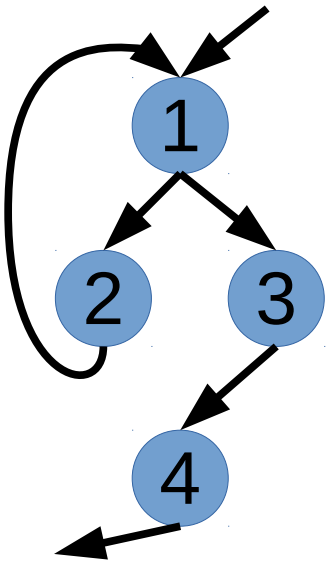
Control-Flow Integrity (CFI)



- Statically construct Control-Flow Graph
  - Find set of allowed targets for each location
- Online set check



Control-Flow Integrity (CFI)



- Statically construct Control-Flow Graph
  - Find set of allowed targets for each location
- Online set check

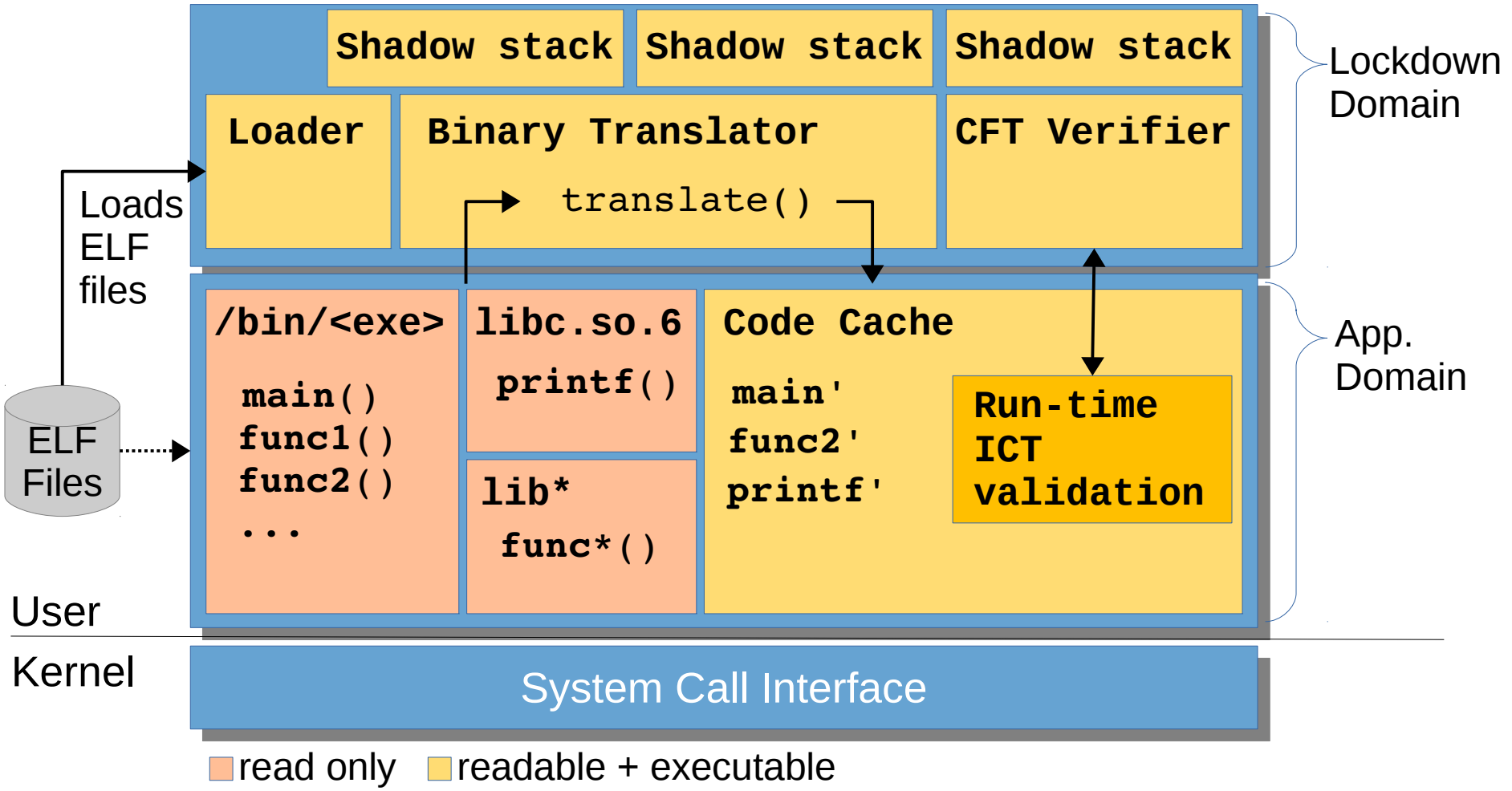
Attacker may write to memory, code pointers verified if used

Fine-grained CFI for binaries

- Fine-grained CFI relies on source code
- Coarse-grained CFI is imprecise

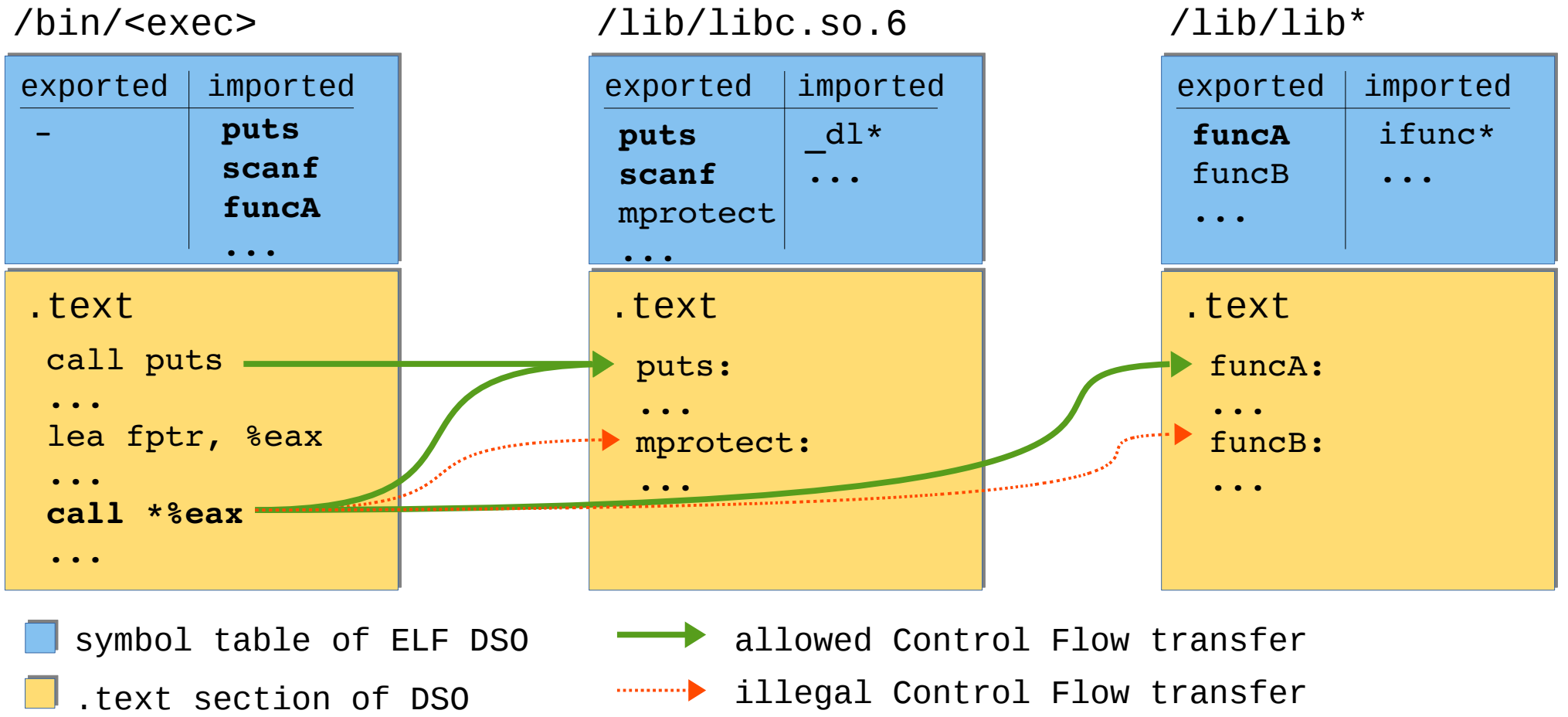
- Goal: enforce fine-grained CFI for binaries
  - Support legacy, binary code
  - Support modularity (libraries)
  - Leverage precise, dynamic analysis
  - Low performance overhead

Lockdown design



Dynamic CFI analysis

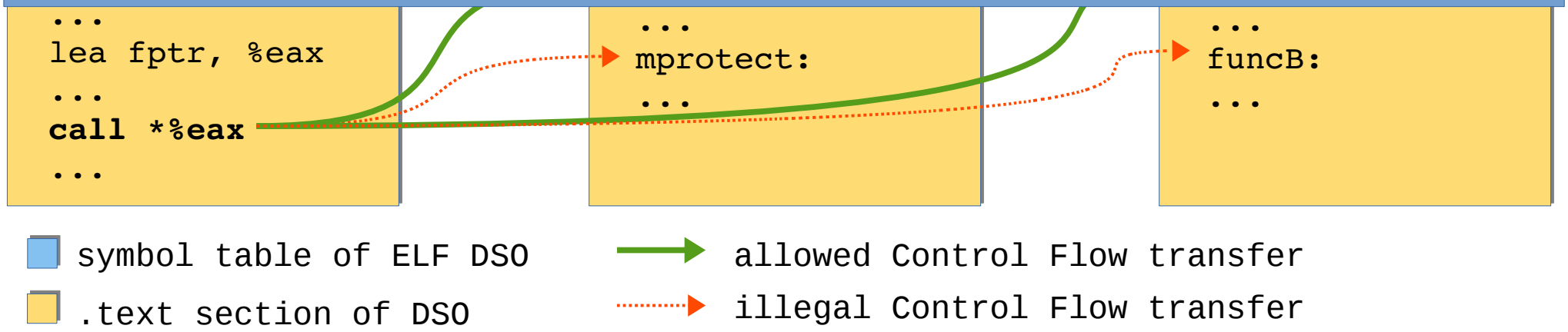
- Leverage program's modularity through loader



Dynamic CFI analysis

- Leverage program's modularity through loader

Modularity increases precision.
No source needed.
Leverage context of transfers.



Lockdown CFI rules

- Return instructions must return to the caller
  - Precise due to shadow stack
- Call instructions must target valid functions
  - Imported in the current module (context)
- Jump instructions must target valid instructions inside the current symbol (or functions)

Performance: Apache 2.2

- 15,000,000 requests
- 56 kB HTML file, 1054 kB image
- Apache 2.2 runs under default configuration

Configuration	Small file	Image	Combined
Single threaded	30.41%	1.94%	7.87%
Concurrent	6.27%	1.09%	1.83%
Concurrent with keep-alive	15.80%	3.00%	4.36%

Security evaluation

- CVE 2013-2028 compromises nginx
  - Both ROP (ret) or COP (icall) exploitation possible

	Length	RET	CALL/JMP/ SYS
ROP attack	30	7	0
COP attack	30	0 (487*)	99

* reachable, but protected by shadow stack

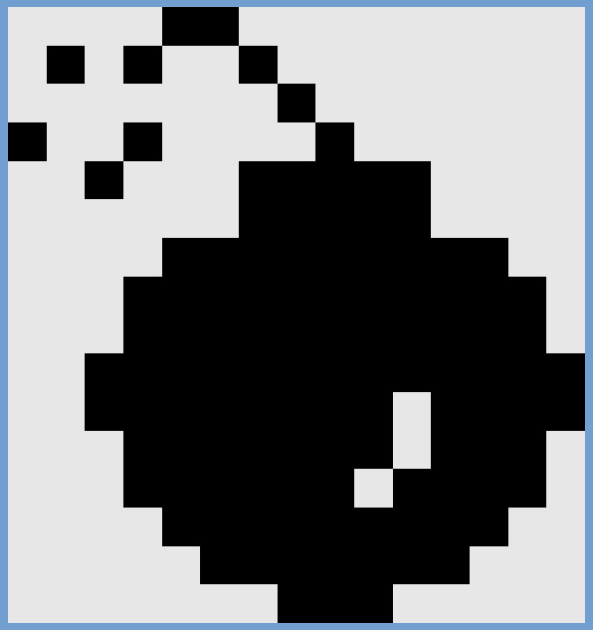
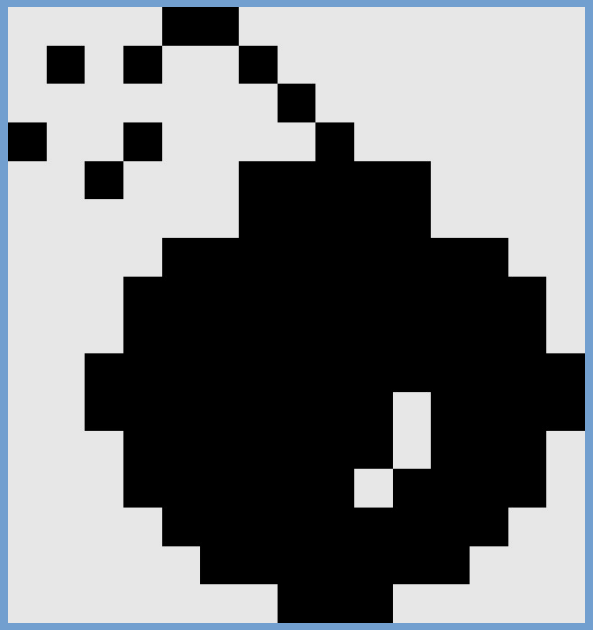
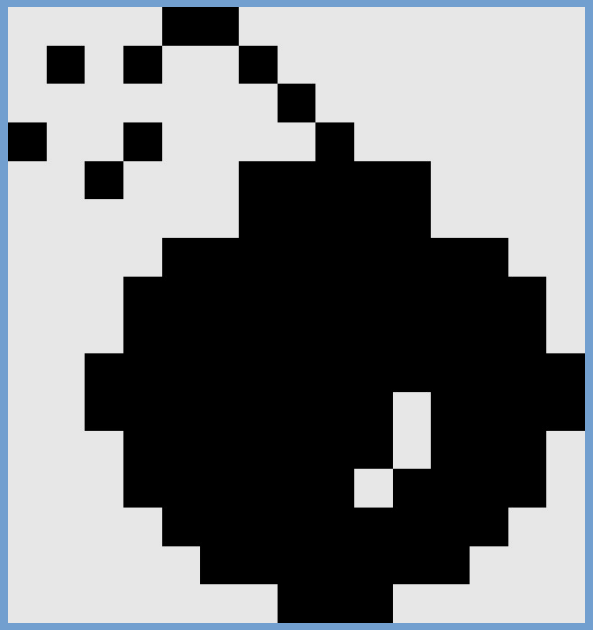
Necessity of shadow stack

- Defenses without stack integrity are broken
  - Loop through two calls to the same function
  - Choose any caller as return location
- Lockdown enforces a protected shadow stack
  - Attacker restricted to arbitrary targets ***on*** the stack
  - Each target can only be called once, in sequence

Conclusion

Conclusion

- Protect in the presence of bugs
- Supports legacy and binary code
- Control-flow hijack protection
  - Shadow stack, dynamic CFI, and locality
  - System call policy as secondary protection
- Reasonably low overhead

Thank you! Questions?

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Performance: SPEC CPU2006